ANY-1 Reference Guide

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# Programming Model

## **Registers**

### Overview

The ANY-1 is a vector machine. The ISA is a 64-register machine with a unified register file for integer, floating-point, decimal-floating-point or posit arithmetic. There are many control and special/status (CSR) registers which hold an assortment of specific values relevant to processing.

General Purpose Registers (x0 to x63) / Scalar Registers

The register usage convention probably has more to do with software than hardware. Excepting a few special cases, the registers are general purpose in nature. Registers may hold integer, floating-point, decimal floating-point or posit values.

x0 always has the value zero. Register x63 is a read only alias of the instruction pointer register. Registers x61 and x62 are used for stack references and subject to stack bounds checking.

|  |  |  |
| --- | --- | --- |
| Register | Description / Suggested Usage | Saver |
| x0 | always reads as zero (hardware) |  |
| x2 | constant building / temporary (cb) |  |
| x3-x9 | temporaries (t0-t6) | caller |
| x10-x19 | register variables (s0-s9) | callee |
| x20-x27 | function arguments (a0-a7) a7/g2 | caller |
| x59 | thread pointer (tp / g1) |  |
| x60 | global data pointer (g0) | callee |
| x61 | base / frame pointer (fp) | callee |
| x62 | current stack pointer (sp) | callee |
|  |  |  |
| x63 | instruction pointer |  |

## Base Registers

Base registers are used as part of the memory management unit of the processing core and are further described in the mmu section of the document.

|  |  |  |
| --- | --- | --- |
| Base Regno | Usage | Selected By |
| b0 to b7 | data | bits 60 to 63 of effective address |
| b8, b9 | reserved | bits 60 to 63 of effective address |
| b10 | Stack | bits 60 to 63 of effective address |
| b11 | I/O | bits 60 to 63 of effective address |
| b12 to b15 | code | bits 60, 63 of instruction pointer |

## Control and Status Registers

Overview

There are numerous special purpose control and status registers in the design. Some registers are present to store variables for performance reasons that would otherwise be stored in main memory.

[U/S/H/M/D]\_CAUSE (0x?006)

This register contains a code indicating the cause of an exception or interrupt. The break handler will examine this code to determine what to do. Only the low order 16 bits are implemented. The high order bits read as zero and are not updateable.

U\_SEMA (CSR 0x000C) Semaphores

This register is available for user semaphores or flag use. Bits in this CSR may be set or cleared with one of the CSRxx instructions. This register has individual bit set / clear capability.

S\_PTA (0x1003)

This register contains the base address of the highest-level page directory for memory management, the paging table depth and the size of the pages mapped. The base address must be page aligned (4kB).

|  |  |  |  |  |
| --- | --- | --- | --- | --- |
| 63 12 | 11 | 10 8 | 7 6 0 | |
| Paging Directory Base Address63..12 | ~ | TD | S | ~ |

|  |  |  |  |  |
| --- | --- | --- | --- | --- |
| TD |  |  | S |  |
| 0 | 1 level lookup |  | 0 | map 16kB pages |
| 1 | 2 level lookup |  | 1 | map 4MB pages |
| 2 | 3 level lookup |  |  |  |
| 3 | 4 level lookup |  |  |  |
| 4 to 7 | reserved |  |  |  |

S\_TID (CSR 0x1010)

This CSR register is reserved for use to contain the task id for the currently running task.

S\_ASID – (CSR 0x101F)

This register contains the address space identifier (ASID) or memory map index (MMI). The ASID is used in this design to select (index into) a memory map in the paging tables.

S\_KEYS – (CSR 0x1020 to 0x1022)

These registers contain a collection of keys associated with the process for the memory system. Each key is twenty bits in size. Each register contains three keys for a total of nine keys. All three registers are searched in parallel for keys matching the one associated with the memory page. Keyed memory enhances the security and reliability of the system.

|  |  |  |  |
| --- | --- | --- | --- |
| 63 60 | 59 40 | 39 20 | 19 0 |
| ~4 | key3 | key2 | key1 |

### M\_BADADDR (CSR 0x3007)

This register contains the effective address for a load / store operation that caused a memory management exception or a bus error. Note that the address of the instruction causing the exception is available in the XL register.

### M\_BAD\_INSTR (CSR 0x300B)

This register contains a copy of the exceptioned instruction.

M\_TVEC (0x3030 to 0x3033)

These registers are an alias for the D\_TVEC registers 0x4030 to 0x4033 which allows access to the TVEC registers from machine mode. They contain the address of the exception handling routine for a given operating level. The lower bits of the exception address are determined from the operating level. TVEC[1] to TVEC[4] are used by the REX instruction.

D\_TVEC (0x4030 to 0x4034)

These registers contain the address of the exception handling routine for a given operating level. TVEC[4] (0x4034) is used directly by hardware to form an address of the debug routine. The lower eight bits of TVEC[4] are not used. The lower bits of the exception address are determined from the operating level. TVEC[1] to TVEC[4] are used by the REX instruction.

### M\_PM\_STACK (0x3040)

This register contains an eight-entry operating mode and interrupt mask stack. When an exception or interrupt occurs, this register is shifted to the left by four bits and the low order bits are set according to the exception mode, when an RTI instruction is executed this register is shifted to the right by four bits. On RTI the last stack entry is set to $9 masking all interrupts on stack underflow. The low order four bits represent the current operating mode and interrupt mask. Only the low order 32 bits of the register are implemented.

### M\_STATUS (0x3044)

This register contains the interrupt mask, operating mode, and privilege level.

|  |  |  |
| --- | --- | --- |
| Bitno | Field | Description |
| 0 to 3 | IM | active interrupt mask level |
| 4 to 5 | ~ | reserved |
| 6 to 13 | PL | privilege level |
| 14 to 19 | RS | register set selection – general purpose registers, this also controls which bounds register set is viewable in the CSRs. |
| 20 | SX | software exception – typically set by a throw() operation and cleared in a catch() handler. |
| 20 to 21 | ~ | reserved |
| 24 to 27 | Thrd | active thread |
| 28 to 31 | IRQ | The level of interrupt that caused the hardware BRK. |
| 32 | VCA | indicates that vector chaining was active prior to an exception |
|  |  |  |
| 40 to 47 | ~ | reserved |
| 48 to 49 | FS | floating point state |
| 50 to 51 | XS | additional core extension state |
| 55 | MPRV | memory privilege: This bit when true (1) causes memory operations to use the first stack privilege level when evaluating privilege and protection rules. (Bits 0 to 31 in the pm\_stack reg). |
| 56 to 60 | VM | These bits control virtual memory options. Note that multiple options may be present at the same time. At reset all the bits are set to zero. |
| 63 | SD |  |

#### VM5

|  |  |  |
| --- | --- | --- |
| Bit | Indicates |  |
| 0 | 1 = single bound |  |
| 1 | 1 = separate program and data bounds |  |
| 2 | 1 = lot protection system |  |
| 3 | 1 = simplified paged unit |  |
| 4 | 1 = paging unit |  |

### M\_EIP (0x3048)

This register contains the interrupt or exception instruction pointer register.

### D\_TIME (0x?FE0)

The TIME register corresponds to the wall clock real time. This register can be used to compute the current time based on a known reference point. The register value will typically be a fixed number of seconds offset from the real wall clock time. The lower 32 bits of the register are driven by the tm\_clk\_i clock time base input which is independent of the cpu clock. The tm\_clk\_i input is a fixed frequency used for timing that cannot be less than 10MHz. The low order 32 bits represent the fraction of one second. The upper 32 bits represent seconds passed. For example, if the tm\_clk\_i frequency is 100MHz the low order 32 bits should count from 0 to 99,999,999 then cycle back to 0 again. When the low order 32 bits cycle back to 0 again, the upper 32 bits of the register is incremented. The upper 32 bits of the register represent the number of seconds passed since an arbitrary point in the past.

Note that this register has a fixed time basis, unlike the TICK register whose frequency may vary with the cpu clock. The cpu clock input may vary in frequency to allow for performance and power adjustments.

## Operating Modes

The core has five operating modes. The highest operating mode is operating mode four which is called the debug operating mode. Debug operating mode has complete access to the machine including special registers and features reserved for debug. Other operating levels may have more restricted access. When an interrupt occurs, the operating mode is set to the debug mode. The core vectors to an address depending on the current operating mode. When not operating at user mode addresses are not subjected to translation and the virtual address and physical address are the same.

|  |  |
| --- | --- |
| Operating Mode | Moniker |
| 0 | user |
| 1 | supervisor |
| 2 | hypervisor |
| 3 | machine |
| 4 | debug |

### Switching Operating Modes

The operating mode is automatically switched to the debug mode when an interrupt occurs. The BRK instruction may be used to switch operating modes. The REX instruction may also be used by a handler to switch the operating mode to a lower mode. One of the exception return instructions (RTD, RTE) will switch the operating level back to what it was prior to the interrupt or exception.

# Exceptions

## External Interrupts

There is little difference between an externally generated exception and an internally generated one. An externally caused exception will force a BRK instruction into the instruction stream. The BRK instruction contains a cause code identifying the external interrupt source.

## Polling for Interrupts

To support code that needs to run with interrupts disabled an interrupt polling instruction (PFI) is provided in the instruction set. For instance, the system could be running a high priority task with interrupts disabled. There may be sections of code where it is possible to process an interrupt however. In some code environments, it is not enough to disable and enable interrupts around critical code. The code must be effectively run with interrupt disabled all the time. This makes it necessary to poll for interrupts in software. For instance, stack prologue code may cause false pointer matches for the garbage collector because stack space is allocated before the contents are defined. If the GC scan occurs on this allocated but undefined area of memory, there could be false matches.

## Effect on Machine Status

The operating mode is always switched to the debug mode on exception. It is up to the debug mode code to redirect the exception to a lower operating mode when desired. Further exceptions at the same or lower interrupt level are disabled automatically. Debug mode code must enable interrupts at some point.

## Exception Stack

The current register set, operating mode and interrupt enable bits are pushed onto an internal stack when an exception occurs. This stack is only eight entries deep as that is the maximum amount of nesting that can occur. Further nesting of exceptions can be achieved by saving the state contained in the exception registers.

## Exception Vectoring

Exceptions are handled through a vector table. The vector table has six entries, one for each operating level the core may be running at. The location of the vector table is determined by TVEC[5]. If the core is operating at mode three for instance and an interrupt occurs vector table address number three is used for the interrupt handler. Note that the interrupt automatically switches the core to operating mode five. An exception handler at the machine level may redirect exceptions to a lower level handler identified in one of the vector registers. More specific exception information is supplied in the cause register.

|  |  |  |
| --- | --- | --- |
| Operating Level | Address (If TVEC[5] contains $F…FC0000) |  |
| 0 | $F…FC0000 | Handler for operating level zero |
| 1 | $F…FC0020 |  |
| 2 | $F…FC0040 |  |
| 3 | $F…FC0060 |  |
| 4 | $F…FC0080 |  |
|  |  |  |

## Reset

The core begins executing instructions at address $F…FD0000. All registers are in an undefined state. Register set #0 is selected.

## Precision

Exceptions in ANY1 are precise. They are processed according to program order of the instructions. If an exception occurs during the execution of an instruction, then an exception field is set in the reorder buffer. The exception is processed when the instruction commits which happens in program order. If the instruction was executed in a speculative fashion, then no exception processing will be invoked unless the instruction makes it to the commit stage.

## Exception Cause Codes

The following table outlines the cause code for a given purpose. These codes are specific to ANY1. Under the HW column an ‘x’ indicates that the exception is internally generated by the processor; the cause code is hard-wired to that use. An ‘e’ indicates an externally generated interrupt, the usage may vary depending on the system.

|  |  |  |  |  |
| --- | --- | --- | --- | --- |
| Cause Code |  | HW | Description |  |
| 0 |  |  | no exception |  |
| 1 | IBE | x | instruction bus error |  |
| 2 | EXF | x | Executable fault |  |
| 4 | TLB | x | tlb miss |  |
|  |  |  | FMTK Scheduler |  |
| 128 |  | e |  |  |
| 129 | KRST | e | Keyboard reset interrupt |  |
| 130 | MSI | e | Millisecond Interrupt |  |
| 131 | TICK | e |  |  |
| 156 | KBD | e | Keyboard interrupt |  |
| 157 | GCS | e | Garbage collect stop |  |
| 158 | GC | e | Garbage collect |  |
| 159 | TSI | e | FMTK Time Slice Interrupt |  |
| 3 |  |  | Control-C pressed |  |
| 20 |  |  | Control-T pressed |  |
| 26 |  |  | Control-Z pressed |  |
|  |  |  |  |  |
| 32 | SSM | x | single step |  |
| 33 | DBG | x | debug exception |  |
| 34 | TGT | x | call target exception |  |
| 35 | MEM | x | memory fault |  |
| 36 | IADR | x | bad instruction address |  |
| 37 | UNIMP | x | unimplemented instruction |  |
| 38 | FLT | x | floating point exception |  |
| 39 | CHK | x | bounds check exception |  |
| 40 | DBZ | x | divide by zero |  |
| 41 | OFL | x | overflow |  |
|  |  |  |  |  |
| 47 |  |  |  |  |
| 48 | ALN | x | data alignment |  |
| 49 | KEY | x | memory key fault |  |
| 50 | DWF | x | Data write fault |  |
| 51 | DRF | x | data read fault |  |
| 52 | SGB | x | segment bounds violation |  |
| 53 | PRIV | x | privilege level violation |  |
| 54 | CMT | x | commit timeout |  |
| 55 | BT | x | branch target |  |
| 56 | STK | x | stack fault |  |
| 57 | CPF | x | code page fault |  |
| 58 | DPF | x | data page fault |  |
| 60 | DBE | x | data bus error |  |
| 61 | PMA | x | physical memory attributes check fail |  |
| 62 | NMI | x | Non-maskable interrupt |  |
|  |  |  |  |  |
| 225 | FPX\_IOP | x | Floating point invalid operation |  |
| 226 | FPX\_DBZ | x | Floating point divide by zero |  |
| 227 | FPX\_OVER | x | floating point overflow |  |
| 228 | FPX\_UNDER | x | floating point underflow |  |
| 229 | FPX\_INEXACT | x | floating point inexact |  |
| 231 | FPX\_SWT | x | floating point software triggered |  |
| 239 |  |  | Software exception handling |  |
| 240 | SYS |  | Call operating system (FMTK) |  |
| 241 |  |  | FMTK Schedule interrupt |  |
| 242 | TMR | x | system timer interrupt |  |
| 243 | GCI | x | garbage collect interrupt |  |
| 253 | RST | x | reset |  |
| 254 | NMI | x | non-maskable interrupt |  |
| 255 | PFI |  | reserved for poll-for-interrupt instruction |  |

### DBG

A debug exception occurs if there is a match between a data or instruction address and an address in one of the debug address registers.

### IADR

This exception is currently not implemented but reserved for the purpose of identifying bad instruction addresses. If the two least significant bits of the instruction address are non-zero then this exception will occur.

### UNIMP

This exception occurs if an instruction is encountered that is not supported by the processor. It may also occur if there is an attempt to use an instruction in a mode that does not support it.

### OFL

If an arithmetic operation overflows (multiply, add, or shift) and the overflow exception is enabled in the arithmetic exception enable register then an OFL exception will be triggered.

### KEY

This fault will occur if an attempt is made to access memory for which the app does not have the key.

### FLT

A floating-point exception is triggered if an exceptional condition occurs in the floating-point unit and the exception is enabled. Please see the section on floating-point for more details.

### DRF, DWF, EXF

Data read fault, data write fault, and execute fault are exceptions that are returned by the memory management unit when an attempt is made to access memory for which the corresponding access type is not allowed. For instance, if the memory page is marked as non-executable an attempt is made to load the instruction cache from the page then an execute fault EXF exception will occur.

### CPF, DPF

The code page fault and data page fault exceptions are activated by the mmu if the page is not present in memory. Access may be allowed but simply unavailable. These faults are not currently implemented.

### PRIV

Some instructions and CSR registers are legal to use only at a higher operating level. If an attempt is made to use the privileged instruction by a lower operating level, then a privilege violation exception may occur. For instance, attempting to use RTI instruction from user operating level.

### STK

If the value loaded into one of the stack pointer registers (the stack pointer sp or frame pointer fp) is outside of the bounds defined by the stack bounds registers, then a stack fault exception will be triggered.

### DBE

A timeout signal is typically wired to the err\_i input of the core and if the data memory does not respond with an ack\_i signal fast enough an error will be triggered. This will happen most often when the core is attempting to access an unimplemented memory area for which no ack signal is generated. When the err\_i input is activated during a data fetch, an exception is flagged in a result register for the instruction. The core will process the exception when the instruction commits. If the instruction does not commit (it could be a speculated load instruction) then the exception will not be processed.

### PMA

The addressed memory did not pass the physical memory attributes testing. For example a write operation attempted to a ROM address space.

### IBE

A timeout signal is typically wired to the err\_i input of the core and if the instruction memory does not respond with an ack\_i signal fast enough and error will be triggered. This will happen most often when the core is attempting to access an unimplemented memory area for which no ack signal is generated. When the err\_i input is activated during an instruction fetch, a breakpoint instruction is loaded into the cache at the address of the error.

### NMI

Non-maskable interrupt.

### BT

The core will generate the BT (branch target) exception if a branch instruction points back to itself. Branch instructions in this sense include jump (JMP) and call (CALL) instructions.

# Memory Management Unit - MMU

## Introduction

Many systems can benefit from the provision of virtual memory management. Virtual memory may be used to protect the address space of one app from another. Virtual memory can enhance the reliability and security of a system.

The simplified system MMU provides minimalistic base and bound and paging capabilities for a small to mid size system. There are two options available for paging, a simple page map ram, and a software managed TLB. The page mapping ram is not suitable for larger systems as the paging tables would be too large. Base bound and paging are applied only to user mode apps. In other operating modes the system sees a flat address space with no restrictions on access. Base address generation is applied to virtual addresses first to generate a linear address which is then mapped using a paged mapping system. Access rights are governed by the base register since all pages in the based on the same address are likely to require the same access. Support for access rights is optional if it is desired to reduce the hardware cost. To simplify hardware there are no bound registers. Bounds are determined by what memory is mapped into the base address area.

Associated with each memory page and stored in its own table is a memory key. The memory key is matched against the keyset in CSR registers. Access to the memory page is allowed only if one of the keys in the keyset matches the memory key, or if the page is marked generally accessible with the special key of zero. Memory may be shared between apps that share the same memory key.

## Base Registers

The upper address bits of a virtual or effective address are not used for addressing memory and are available to select base register. The MMU includes 16 base registers. The base register in use is selected by the upper nybble of the virtual address. If the program address has all ones in bits 24 to 63 then base addressing is bypassed. This provides a shared program area containing the BIOS and OS code.

|  |  |  |
| --- | --- | --- |
| Base Regno | Usage | Selected By |
| 0 to 7 | data | bits 60 to 63 of effective address |
| 8, 9 | reserved | bits 60 to 63 of effective address |
| 10 | Stack | bits 60 to 63 of effective address |
| 11 | I/O | bits 60 to 63 of effective address |
| 12 to 15 | code | bits 60, 63 of instruction pointer |

### Base Register Format

|  |  |
| --- | --- |
| 63 4 | 3 0 |
| Base Address60 | RWX |

The low order four bits of the base register are reserved for access rights bits. Supporting memory access rights is optional.

R: 1 = segment readable

W: 1 = segment writeable

X: 1 = segment executable

### Base Register Access

Base registers may be read and altered using the MVBASE instruction. The MVBASE instruction works in an indirect fashion as described in the text.

## Linear Address Generation

The base address value contained in the upper 60 bits of a base register is shifted left 14 bits before being added to the virtual address. This gives potentially a 74-bit address space.

Note there is no limit or bound register. Access is limited by what is mapped into the segment. Unmapped page addresses result in a page fault exception. Pages that are inaccessible use the reserved page number of all ones.

## The Page Map

The page directly maps virtual address pages to physical ones. The page map is a dedicated memory internal to the processing core accessible with the custom ‘mvmap’ instruction. It is similar in operation to a TLB but is much simpler. TLB’s cache address translations and create TLB miss exceptions. Page walks of mapping tables are required to update the TLB on a miss. There are no exceptions associated with the page mapping table.

In addition to based addresses, memory is divided up into 16kB pages which are mapped. There are 32 memory maps available. A memory map represents an address space; a five-bit address space identifier is in use. Address spaces will need to be shared if more than 32 apps are running in the system. The desire is to keep the mapping tables small so they may fit into a small number of standard memory blocks. For instance, for the sample system there are 4096 pages required to map the 256MB address space. Any individual app is limited to maximum of 64MB (one quarter of the memory available). The virtual page number is used to lookup the physical page in the page mapping table. Addresses with the top eight bits set are not mapped to allow access to the system ROM. Pages that are inaccessible use the reserved physical page number of all ones.

The page mapping table is indexed by the ASID and the virtual page number to determine the physical page. The ‘mvmap’ instruction uses Rs1 to contain a mapping table index. Bits 16 to 20 of Rs1 are the ASID, bits 0 to 15 of Rs1 are used for the virtual page number. It is expected that the virtual page number is a small number, in this case 12 bits. Rs2 contains the new value of the physical page. The current value of the physical page is placed in Rd when the instruction executes.

### Page Map Entry Layout

|  |  |  |  |  |  |  |  |  |  |  |
| --- | --- | --- | --- | --- | --- | --- | --- | --- | --- | --- |
| 41 36 | 35 24 | 23 | 22 | 21 | 20 | 19 | 18 | 17 | 16 | 15 0 |
| ASID5 | Virtual Page | G | D | A | U | C | R | W | X | Physical Page |

|  |  |  |  |  |  |  |  |  |  |  |
| --- | --- | --- | --- | --- | --- | --- | --- | --- | --- | --- |
| 0 | 0 |  |  |  |  |  |  |  |  | 10 |
| 1 |  |  |  |  |  |  |  |  | 11 |
| … |  |  |  |  |  |  |  |  |  |
| 4094 |  |  |  |  |  |  |  |  | 18 |
| 4095 |  |  |  |  |  |  |  |  | 19 |
| 1 | 0 |  |  |  |  |  |  |  |  |  |
| 1 |  |  |  |  |  |  |  |  |  |
| … |  |  |  |  |  |  |  |  |  |
| 4094 |  |  |  |  |  |  |  |  |  |
| 4095 |  |  |  |  |  |  |  |  |  |
| … 30 more address spaces | |  |  |  |  |  |  |  |  |  |

The low order 14 bits of an address pass through both linear address generation and paging unchanged.

### The 16kB Page

Many memory systems use a 4kB page size. A 16kB page size is used here mainly to restrict the number of page entries in the page map table. A smaller page size would result in too many pages of memory to support multiple tasks. Even given a 16kB page size there are still 4096 pages of memory available in a map.

*The author was tempted to divide the page mapping table into several different regions capable of mapping different amounts of the address space (small, medium, and large areas). This potentially could allow more memory maps to be present while at the same time not increasing the page table size. However, it would add extra complexity to the memory system which is currently simpler in nature.*

### The MVMAP Instruction

The memory mapping table is managed with a dedicated instruction - [MVMAP](#_MVMAP_–_Move). MVMAP allows high-speed access to the mapping table.

*While the memory mapping table could have been managed with CSRs or possibly be mapped into the main memory space, the author feels that having a dedicated instruction makes the software managing the tables simpler and cleaner.*

Rs1:

|  |  |  |
| --- | --- | --- |
| 63 20 | 20 16 | 15 0 |
| Unused - should be zero | ASID5 | Virtual page number 16 bits max |

## TLB – Translation Lookaside Buffer

### Overview

The page map is limited in the translations it can perform because of its size. The solution to allowing more memory to be mapped is to use main memory to store the translations tables, then cache address translations in a translation look-aside buffer or TLB. This is sometimes also called an address translation cache ATC. The TLB offers a means of address virtualization and memory protection. A TLB works by caching address mappings between a real physical address and a virtual address used by software. The TLB deals with memory organized as pages. Typically, software manages a paging table whose entries are loaded into the TLB as translations are required.

### Size / Organization

The TLB has 1024 entries per set. The size was chosen as it is the size of one block ram for 32-bit data in the FPGA. This is quite a large TLB. Many systems use smaller TLBs. There is not really a need for such a large one, however it is available.

The TLB is organized as a four-way set associative cache.

### What is Translated

The TLB processes all user mode addresses including both instruction and data addresses. It is known as a *unified* TLB. Addresses in other modes of operation are not translated. Additionally, addresses with the top forty bits set are not translated to allow access to the BIOS and system rom.

### Page Size

Because the TLB caches address translations it can get away with a much smaller page size than the page map can for a larger memory system. 4kB is a common size for many systems. In this case the TLB uses 16kB pages to match the size of pages for keyed memory and segmentation. For a 512MB system (the size of the memory in the test system) there are 32,768 16kB pages.

### Management

The TLB unit is a software managed TLB. When a translation miss occurs, an exception is generated to allow software to update the TLB. It is left up to software to decide how to update the TLB. There may be a set of hierarchal page tables in memory, or there could be a hash table used to store translations.

The TLB is updated using the TLBRW instruction which both reads and writes the TLB. More descriptive text is present at the [TLBRW](#_TLBRW_–_Read) instruction description.

### Flushing the TLB

The TLB maintains the address space (ASID) associated with a virtual address. This allows the TLB translations to be used without having to flush old translations from the TLB during a task switch.

#### Global Bit

In addition to the ASID the TLB entries contain a bit that indicates that the translation is a global translation and should be present in every address space.

## PAM – Page Allocation Map

### Overview

Memory is organized into 32,768 16kB pages.

The PAM is a software structure made up of 32,768 bit-pairs stored in memory. There is a bit pair for each possible physical memory page. The PAM is used by software to manage the allocation of physical pages of memory.

### Memory Usage

Total memory used by the PAM is 8kB.

### Organization

The PAM is organized as a string of bit-pairs, one pair for each physical memory page. Bit pairs are used rather than single bits to mark allocated pages as it is convenient to also mark runs of pages. Marking runs of pages using bit-pairs makes it possible to free the pages of a previous allocation.

|  |  |
| --- | --- |
| Bit-Pair Value | Meaning |
| 0 | Page of memory is free, available for use. |
| 1 | reserved |
| 2 | Page is allocated, end of run of pages |
| 3 | Page is allocated |

## PMA - Physical Memory Attributes Checker

### Overview

The physical memory attributes checker is a hardware module that ensures that memory is being accessed correctly according to its physical attributes.

Physical memory attributes are stored in an eight-entry table. This table includes the address range the attributes apply to and the attributes themselves. Address ranges are resolved only to bit four of the address. Meaning the granularity of the check is 16 bytes.

Most of the entries in the table are hard-coded and configured when the system is built.

Physical memory attributes checking is applied in all operating modes.

### Register Description

|  |  |  |  |
| --- | --- | --- | --- |
| Regno | Bits |  |  |
| 00 | 64 | LB0 | lower bound - address bits 4 to 67 of the physical address range |
| 08 | 64 | UB0 | upper bound - address bits 4 to 67 of the physical address range |
| 10 | 16 | AT0 | memory attributes |
| 18 | ~ | ~ | reserved |
| … | … | … | 6 more register sets |
| E0 | 64 | LB7 | lower bound - address bits 4 to 67 of the physical address range |
| E8 | 64 | UB7 | upper bound - address bits 4 to 67 of the physical address range |
| F0 | 16 | AT7 | memory attributes |
| F8 | ~ | ~ | reserved |

### Attributes

|  |  |  |
| --- | --- | --- |
| Bitno |  |  |
| 0 | X | may contain executable code |
| 1 | W | may be written to |
| 2 | R | may be read |
| 3 | C | may be cached |
| 4-6 | G | granularity   |  |  | | --- | --- | | G |  | | 0 | byte accessible | | 1 | wyde accessible | | 2 | tetra accessible | | 3 | octa accessible | | 4 to 7 | reserved | |
| 7 | ~ | reserved |
| 8-15 | T | device type (rom, dram, eeprom, I/O, etc) |

## Key Cache

### Overview

Associated with each page of memory is a memory key. To access a page of memory the memory key must match with one of the keys in the applications keyset. The keyset is maintained in the keys CSRs. The key size of 20 bits is a minimum size recommended for security purposes.

The key associated with each memory page is stored in a table in main memory. Each key occupies a tetra-byte of memory to keep caching simple. So that two memory accesses are not required to access a page of memory this table of keys is cached. When a page of memory is accessed the key cache is accessed in parallel.

The key cache is a direct mapped cache organized as 256 lines of 16 keys. Key values are stored in LUT rams. 256 address tags are stored in LUT ram.

## Card Memory

### Overview

Also present in the memory system is Card memory. The card memory is a telescopic memory which reflects with increasing detail where in the memory system a pointer write has occurred. This is for the benefit of garbage collection systems. Card memory is updated using a write barrier when a pointer value is stored to memory.

### Organization

Memory is divided into 256-byte card memory pages. Each card has a single byte recording whether a pointer store has taken place in the corresponding memory area. To cover a 512MB memory system 2MB card memory is required at the outermost layer. The outer most 2MB card memory layer is itself divided into 4096 256-byte card pages. Note that each byte represents the pointer store status for a 256B region. The 4096B memory is further resolved to single octa indicating if any pointer store has taken place. Thus, for a 512MB memory system a three-level card memory is used.

|  |  |  |
| --- | --- | --- |
| Layer | Resolving Power | |
| 0 | 2 MB | 256B regions |
| 1 | 4 kB | 128kB regions |
| 2 | 8 B | 64 MB regions |

There is only a single card memory in the system, used by all tasks.

### Location

Card memory must be based at physical address zero, extending up to the amount of card memory required. This is so that the address calculation of the memory update may be done with a simple right-shift operation.

### Operation

As a program progresses it writes pointer values to memory using the write barrier. Storing a pointer triggers an update to all the layers of card memory corresponding to the main memory location written. A byte is set in each layer of the card memory system corresponding to the memory location of the pointer store.

The garbage collection system can very quickly determine where pointer stores have occurred and skip over memory that has not been modified.

### Sample Write Barrier

; Milli-code routine for garbage collect write barrier.

; Usable with up to 64-bit memory systems.

; Three level card memory

;

GCWriteBarrier:

STO a0,[a1] ; store the value to memory at a1

SRL a1,a1,#8 ; compute card address

STB x0,[a1] ; clear bit in card memory

SRL a1,a1,#8 ; repeat for each table level

STB x0,[a1]

SRL a1,a1,#8

STB x0,[a1]

;… more stores as needed

JMP ra1

## System Memory Map

There are several components to the system which use tables in memory. These tables are statically allocated at the time the system is built. The table sizes depend on the size of main memory. The card memory table must be located at address zero. So, it is probably best to group the tables together at the low end of memory.

|  |  |  |
| --- | --- | --- |
| Address | Usage |  |
| $00000000 to $001FFFFF | Card Memory (2 MB) |  |
| $00202000 to $00203FFF | PAM (8kB 2 copies) |  |
| $00280000 to $0029FFFF | Key memory (128 kB) |  |
|  |  |  |